



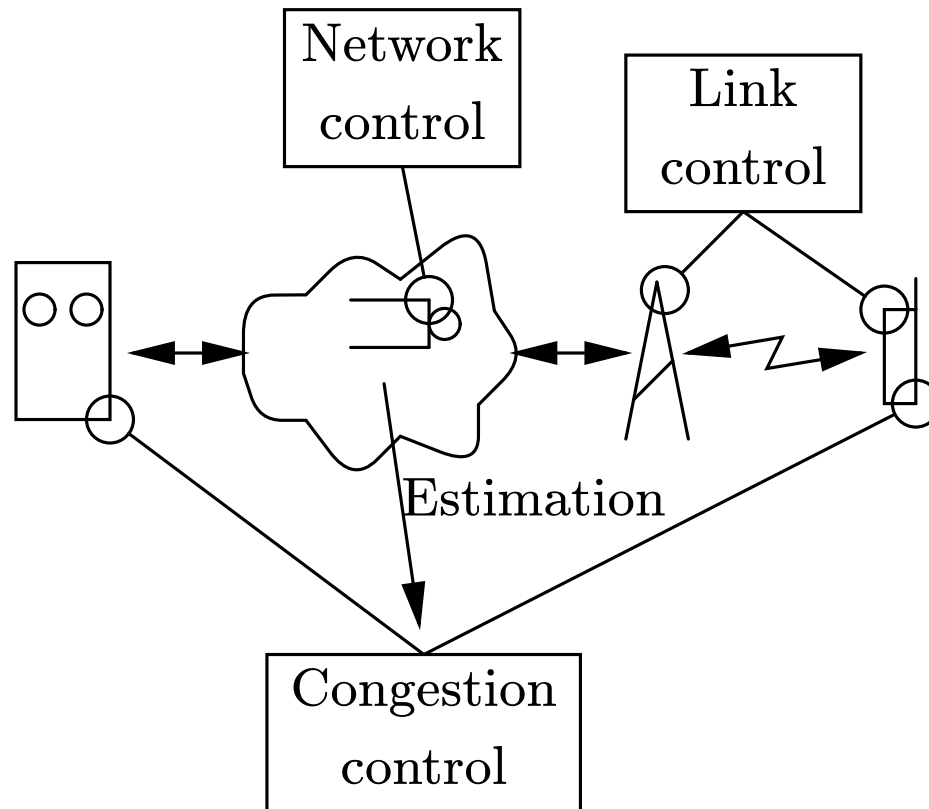
SOME MODELING AND ESTIMATION ISSUES IN CONTROL OF HETEROGENOUS NETWORKS

Krister Jacobsson, Niels Möller,
Karl Henrik Johansson and Håkan Hjalmarsson
Department of Signals, Sensors and Systems, KTH

Automatic Control



OVERVIEW



Focus of our group:

- Model based estimation.
- Making wireless links friendlier to TCP.

MODEL BASED ESTIMATION

Window-based flow-control objective:

$$w \approx b \cdot \text{RTT}$$

- Estimation of round-trip time (RTT).
- Estimation of available bandwidth (b).
- Trade-off between noise reduction and tracking performance.
- Model-based estimation. Systematic way to make that trade-off.

RTT ESTIMATION

Model:

- Piecewise constant “average RTT” x_k .
- Occasional step changes due to rerouting, bottlenecks appearing. . .

$$x_{k+1} = x_k + \delta_k v_k \quad \delta_k \in \{0, 1\}$$

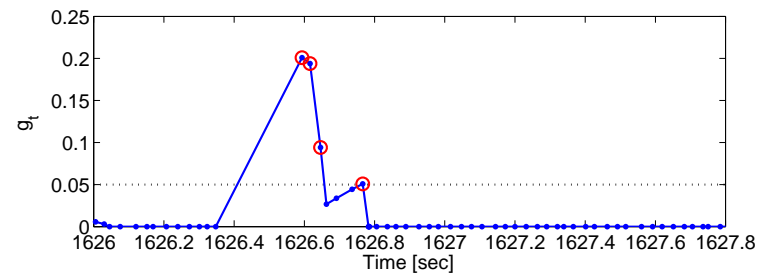
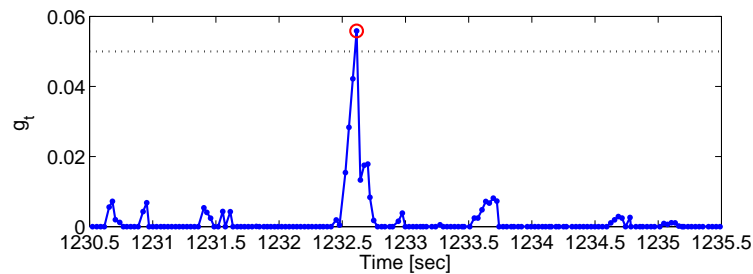
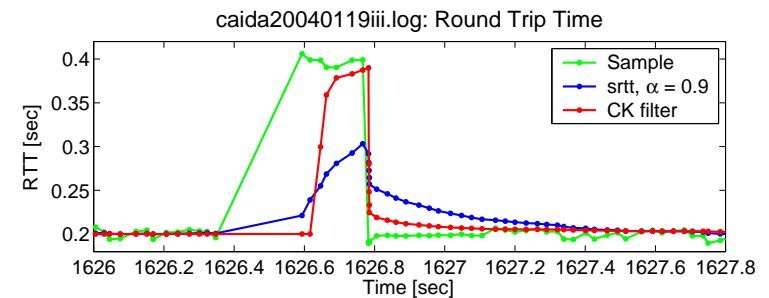
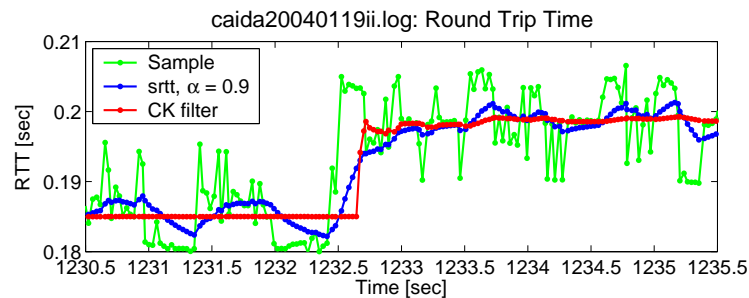
$$y_k = x_k + e_k \quad \text{Measured RTT}$$

Proposed estimator:

- Kalman filter to suppress noise.
- Change detection to track δ_k .

EVALUATION OF RTT ESTIMATORS

Input: RTT samples KTH \leftrightarrow Caida, ≈ 20 hops, interval 30 ms.



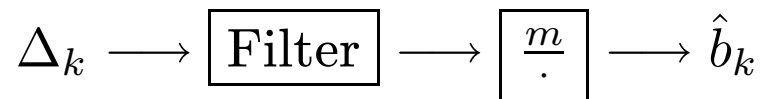
Bottom: Output from the change detection.

BANDWIDTH ESTIMATION

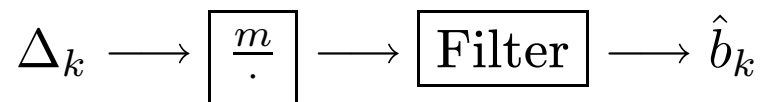
Measurements: ACK inter-arrival times Δ_k .

$$b_N = \frac{Nm}{\sum_k \Delta_k} = \frac{m}{\frac{1}{N} \sum_k \Delta_k} \quad \text{Constant packet size } m$$

Model: $\Delta_k = b + e_k$, zero-mean noise e_k . Use a low-pass filter:



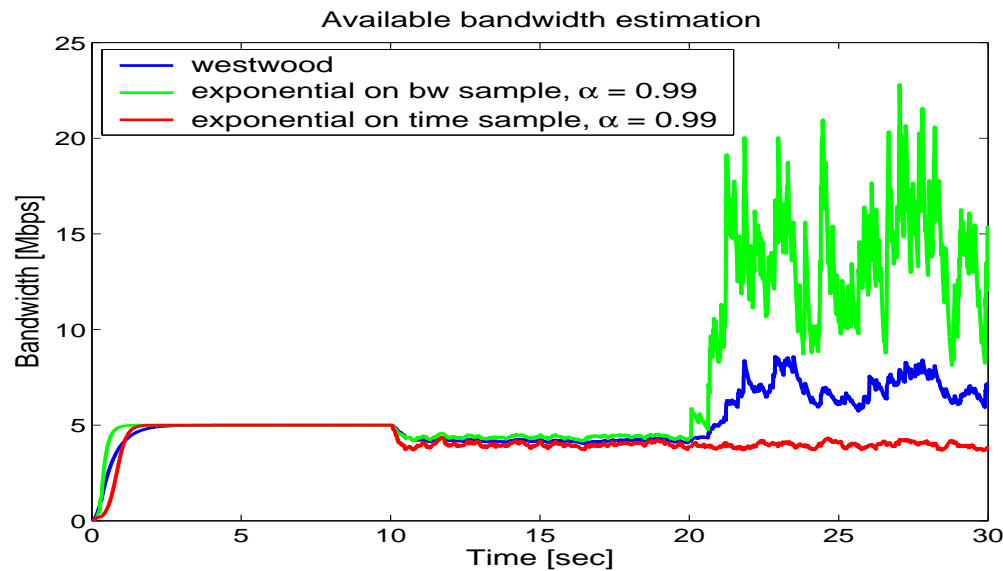
Alternative structure (used in TCP-Westwood):



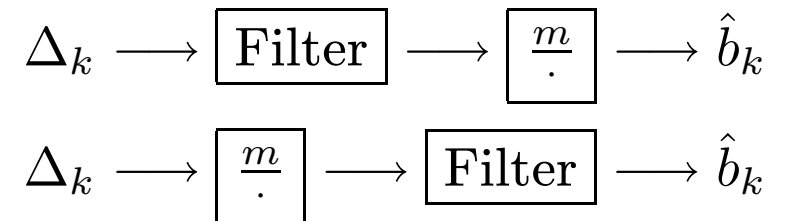
Results in bias *independent* of filter design.

EVALUATION OF BANDWIDTH ESTIMATORS

Input: TCP simulation in ns-2, 5 Mbps bottleneck.



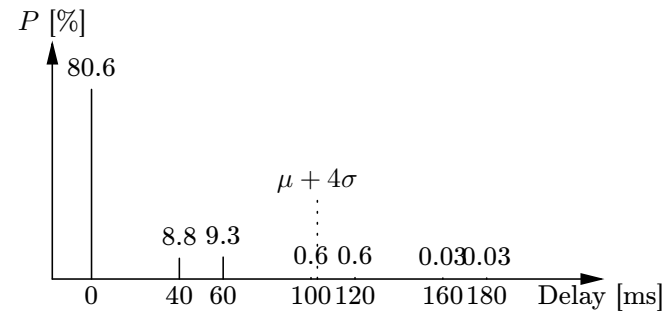
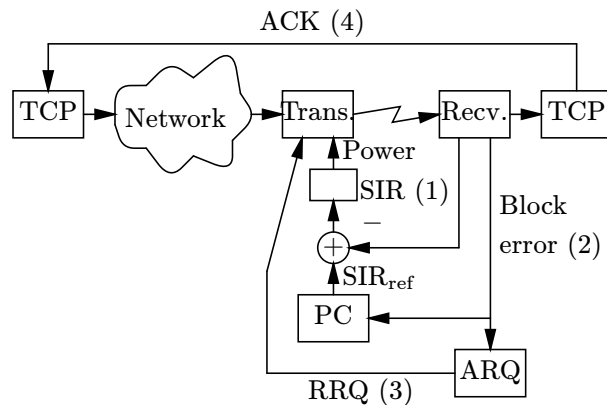
Should filter before the non-linearity.



At 10 ms: 1 Mbps cross-traffic in forward direction.

At 20 ms: 1 Mbps cross-traffic in reverse direction.

INFLUENCE OF WIRELESS LINKS ON TCP



Without link-layer retransmissions: Constant delay, high loss-rate.

With link-layer retransmissions: Random delay, small loss-rate.

Link delay distribution influences TCP. Spurious timeouts.

A MEASURE OF TCP-FRIENDLINESS

Let X be the stochastic link delay.

$$P_{\text{TO}}(X) := P(X > E(X) + 4\sigma(X)) \quad P(\text{Timeout}) \text{ for TCP}$$

- Uniform distribution: $P_{\text{TO}} = 0$.
- Normal distribution: $P_{\text{TO}} \approx 0.006\%$.
- General distribution: $P_{\text{TO}} \leq 6.25\%$.
- Wireless link: $P_{\text{TO}} \approx 0.7\%$.

TWEAKING THE DELAY

Original: $P(X = d_i) = p_i$.

Tweaked: $P(\tilde{X} = d_i + x_i) = p_i$.

$$\min E(\tilde{X})$$

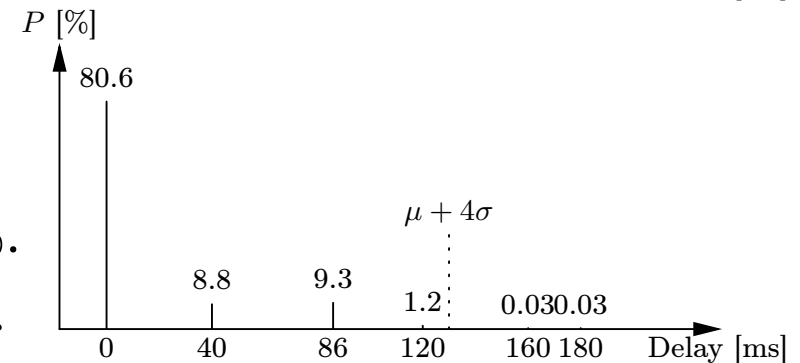
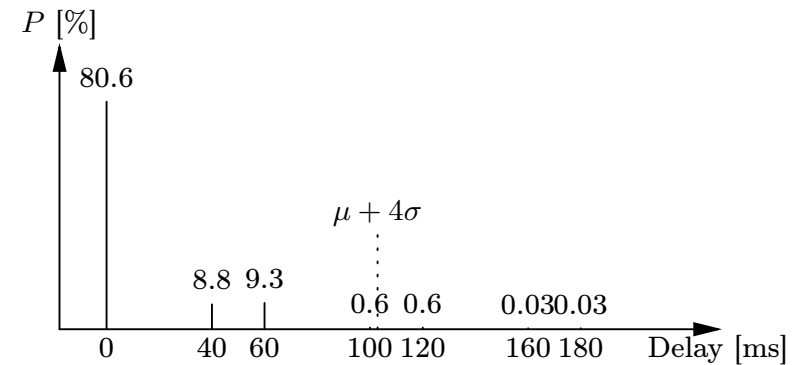
$$P_{\text{TO}}(\tilde{X}) < \epsilon$$

$$x_i \geq 0$$

Decreased P_{TO} , from 0.68% to 0.06%.

Mean delay increased by only 2.5 ms.

Eliminates most spurious timeouts.



CONCLUSIONS

- RTT estimation: Promising model based approach.
- Bandwidth estimation: Average inter-arrival times, not “bandwidth samples”.
- Artificial delays at the link-layer improve TCP performance.
- For wireless links: Use engineering freedom in the link layer.

Vision: Systematic design of network control mechanisms:

- End-to-end congestion control.
- Network-layer control in intermediate routers.
- Link-layer control loops.